

MAT 685: C++ for Mathematicians

Prime pairs arrayed

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Outline

- 1 Prime pairs, yet again
- 2 Euler's totient function
 - A first factoring algorithm
 - A better factoring algorithm
 - Back to the totient
- 3 Counting pairs
- 4 Answering the question
- 5 Summary

Prime pairs, yet
again

Euler's totient
function

A first factoring
algorithm

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Classic problem

- Choose n
- Choose $a, b \in \{1, \dots, n\}$
- Let p_n be probability that $\gcd(a, b) = 1$
- Does $\lim_{n \rightarrow \infty} p_n$ exist?
 - If so, what is its value?

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 - If so, what is its value?

Example

Let $n = 8$.

- Possible outcomes: $\binom{8}{2} = \frac{8!}{2!6!} = 28$
- Relatively prime pairs: $(1, 2), (1, 3), \dots, (1, 8), (2, 3), (2, 5), (2, 7), (3, 4), (3, 5), (3, 7), (3, 8), (4, 5), (4, 7), (6, 7), (7, 8)$
- So $p_8 = 18/28 = 9/14 \approx 64.3\%$

A problem

Still too many numbers! Even the Monte Carlo algorithm takes too long at some point.

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Get Euler's help!

Euler's "totient" function

$$\varphi(n) = |\{m \in \mathbb{N} : 1 \leq m \leq n \text{ and } \gcd(m, n) = 1\}|.$$

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Euler's "totient" function

$$\varphi(n) = |\{m \in \mathbb{N} : 1 \leq m \leq n \text{ and } \gcd(m, n) = 1\}|.$$

Example

n	2	3	4	5	6	7	8	9	10
$\varphi(n)$	1	2	2	4	2	6	4	6	4

n	11	12	13	14	15	16	17
$\varphi(n)$	10	4	12	6	8	8	16

Notice anything?

n	2	3	4	5	6	7	8	9	10
$\varphi(n)$	1	2	2	4	2	6	4	6	4

n	11	12	13	14	15	16	17
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- if $n = p$ is prime...
- if $n = p^k$ is a prime power...
- if $n = ab$ is a relatively prime product...

Notice anything?

n	2	3	4	5	6	7	8	9	10
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- if $n = p$ is prime... $\varphi(p) = p - 1$
- if $n = p^k$ is a prime power...
- if $n = ab$ is a relatively prime product...

Notice anything?

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- if $n = p$ is prime... $\varphi(p) = p - 1$
- if $n = p^k$ is a prime power... $\varphi(p^k) = p^k - p^{k-1}$
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Notice anything?

n	2	3	4	5	6	7	8	9	10
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- if $n = p$ is prime... $\varphi(p) = p - 1$
- if $n = p^k$ is a prime power... $\varphi(p^k) = p^k - p^{k-1}$
- if $n = ab$ is a relatively prime product... $\varphi(ab) = \varphi(a) \varphi(b)$

These properties make sense!

Theorem

If n is prime, then $\varphi(n) = n - 1$.

Proof.

Think about it a moment. □

Theorem

Proof.

□

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If n is prime, then $\varphi(n) = n - 1$.

Proof.

$1, 2, \dots, n - 1$ are all rel. prime to n . □

Theorem

Proof.

□

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If n is prime, then $\varphi(n) = n - 1$.

Proof.

1, 2, ..., $n - 1$ are all rel. prime to n . □

Theorem

If $n = p^k$ is a prime power, then $\varphi(n) = p^k - p^{k-1}$.

Proof.

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These properties make sense!

Theorem

If n is prime, then $\varphi(n) = n - 1$.

Proof.

$1, 2, \dots, n - 1$ are all rel. prime to n . □

Theorem

If $n = p^k$ is a prime power, then $\varphi(n) = p^k - p^{k-1}$.

Proof.

Only $p, 2p, 3p, \dots, (p^{k-1})p$ are not rel. prime to p . □

This one we can only sketch

Theorem

If $n = ab$ is a relatively prime product, then $\varphi(n) = \varphi(a)\varphi(b)$.

Proof.

Think about it a moment.



This one we can only sketch

Theorem

If $n = ab$ is a relatively prime product, then $\varphi(n) = \varphi(a)\varphi(b)$.

Proof.

Products of numbers not rel. prime to a or b are also not rel. prime to p . Numbers not rel. prime to p must also have a common factor with a or b (requires Chinese Remainder Theorem). □

Non-trivial example

Compute $\varphi(100)$

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$$\varphi(ab) = \varphi(a) \varphi(b)$$

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Compute $\varphi(100)$

- $\varphi(100) = \varphi(2^2 \times 5^2)$
- $\varphi(100) = \varphi(2^2) \times \varphi(5^2)$
- $\varphi(100) = (2^2 - 2^1) \times (5^2 - 5^1)$

$$\varphi(ab) = \varphi(a) \varphi(b)$$

$$\varphi(p^k) = p^k - p^{k-1}$$

Non-trivial example

Compute $\varphi(100)$

- $\varphi(100) = \varphi(2^2 \times 5^2)$
- $\varphi(100) = \varphi(2^2) \times \varphi(5^2)$
- $\varphi(100) = (2^2 - 2^1) \times (5^2 - 5^1)$
- $\varphi(100) = 40$

$$\varphi(ab) = \varphi(a) \varphi(b)$$

$$\varphi(p^k) = p^k - p^{k-1}$$

Non-trivial example

Compute $\varphi(100)$

- $\varphi(100) = \varphi(2^2 \times 5^2)$
- $\varphi(100) = \varphi(2^2) \times \varphi(5^2)$
- $\varphi(100) = (2^2 - 2^1) \times (5^2 - 5^1)$
- $\varphi(100) = 40$

$$\varphi(ab) = \varphi(a) \varphi(b)$$

$$\varphi(p^k) = p^k - p^{k-1}$$

Bingo!

The 40 numbers are

01, 03, 07, 09, 11, 13, 17, 19, ..., 91, 93, 97, 99.

(Ten groups of 4.)

How does all this help?

$n = 10$: (1 extra for $(1, 1)$)

$$\begin{aligned}
 p_n &= \frac{1 + 2 \times |\{\text{rel. prime pairs } (i, j), i < j\}|}{|(a, b) : 1 \leq a, b \leq 10|} \\
 &= \frac{1 + 2 \times |\{(1, 2)\} \cup \{(1, 3), (2, 3)\} \cup \dots \cup \{(1, 10), (3, 10), \dots, (9, 10)\}|}{100} \\
 &= \frac{1 + 2 \times (|\{(1, 2)\}| + \dots + |\{(1, 10), (3, 10), \dots, (9, 10)\}|)}{100} \\
 &= \frac{1 + 2 \times (\varphi(2) + \varphi(3) + \dots + \varphi(10))}{100}
 \end{aligned}$$

Prime pairs, yet again

Euler's totient function

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 p_n &= \frac{1 + 2 \times |\{\text{rel. prime pairs } (i, j), i < j\}|}{|(a, b) : 1 \leq a, b \leq 10|} \\
 &= \frac{1 + 2 \times |\{(1, 2)\} \cup \{(1, 3), (2, 3)\} \cup \dots \cup \{(1, 10), (3, 10), \dots, (9, 10)\}|}{100} \\
 &= \frac{1 + 2 \times (|\{(1, 2)\}| + \dots + |\{(1, 10), (3, 10), \dots, (9, 10)\}|)}{100} \\
 &= \frac{1 + 2 \times (\varphi(2) + \varphi(3) + \dots + \varphi(10))}{100}
 \end{aligned}$$

In general,

$$p_n = \frac{1 + 2 \sum_{k=2}^n \varphi(k)}{n^2}$$

(book gives different, but equivalent, formula)

So what?

$$p_n = \frac{1 + 2 \sum_{k=2}^n \varphi(k)}{n^2}$$

given n

let $s = 0$

for each $k \in \{2, \dots, n\}$

add $\varphi(k)$ to s

multiply 2 to s

add 1 to s

divide s by n^2

return s

So what?

$$p_n = \frac{1 + 2 \sum_{k=2}^n \varphi(k)}{n^2}$$

given n
let $s = 0$
for each $k \in \{2, \dots, n\}$
 add $\varphi(k)$ to s
multiply 2 to s
add 1 to s
divide s by n^2
return s

Example ($n = 10$)

start w/s = 0
 $k = 2, 3, \dots, 10$
 $s = 1, 3, 5, 9, 11, 17, 21, 27, 31$
multiply: $s = 62$
add: $s = 63$
divide: $s = 0.63$
return 0.63

p. 88 #5.1

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Factoring n

To find $\varphi(n)$, we need n 's factors

Question

How do we find them?

Factoring n

To find $\varphi(n)$, we need n 's factors

Question

How do we find them?

given n

let L be a list of \sqrt{n} zeroes

for each $i \in \{2, \dots, \sqrt{n}\}$

while $i \mid n$

 increment L_i

 replace n by n/i

return L

Example

$$n = 100$$

$$L = (0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$$

Example

$$n = 100$$

$$L = (0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$$

loop $i = 2$

Example

$$n = 100$$

$$L = (0, 1, 0, 0, 0, 0, 0, 0, 0, 0)$$

loop $i = 2$

$2 \mid 100$, so increment L_2 and replace n by 50

Example

$$n = 100$$

$$L = (0, 2, 0, 0, 0, 0, 0, 0, 0)$$

loop $i = 2$

$2 \mid 100$, so increment L_2 and replace n by 50

$2 \mid 50$, so increment L_2 and replace n by 25

Example

$n = 100$

$L = (0, 2, 0, 0, 0, 0, 0, 0, 0)$

loop $i = 2$

$2 \mid 100$, so increment L_2 and replace n by 50

$2 \mid 50$, so increment L_2 and replace n by 25

$2 \nmid 25$: **while** loop ends

Example

$n = 100$

$L = (0, 2, 0, 0, 0, 0, 0, 0, 0)$

loop $i = 2$

2 | 100, so increment L_2 and replace n by 50

2 | 50, so increment L_2 and replace n by 25

2 † 25: **while** loop ends

loop $i = 3$

3 † 25: **while** loop ends

Example

$n = 100$

$L = (0, 2, 0, 0, 0, 0, 0, 0, 0)$

loop $i = 2$

2 | 100, so increment L_2 and replace n by 50

2 | 50, so increment L_2 and replace n by 25

2 † 25: **while** loop ends

loop $i = 3$

3 † 25: **while** loop ends

loop $i = 4$

4 † 25: **while** loop ends

Example

 $n = 100$ $L = (0, 2, 0, 0, 0, 0, 0, 0, 0)$ loop $i = 2$ $2 \mid 100$, so increment L_2 and replace n by 50 $2 \mid 50$, so increment L_2 and replace n by 25 $2 \nmid 25$: **while** loop endsloop $i = 3$ $3 \nmid 25$: **while** loop endsloop $i = 4$ $4 \nmid 25$: **while** loop endsloop $i = 5$

Example

 $n = 100$ $L = (0, 2, 0, 0, 1, 0, 0, 0, 0)$ loop $i = 2$ $2 \mid 100$, so increment L_2 and replace n by 50 $2 \mid 50$, so increment L_2 and replace n by 25 $2 \nmid 25$: **while** loop endsloop $i = 3$ $3 \nmid 25$: **while** loop endsloop $i = 4$ $4 \nmid 25$: **while** loop endsloop $i = 5$ $5 \mid 25$: increment L_5 and replace n by 5

Example

$n = 100$

$L = (0, 2, 0, 0, 2, 0, 0, 0, 0, 0)$

loop $i = 2$

2 | 100, so increment L_2 and replace n by 50

2 | 50, so increment L_2 and replace n by 25

2 † 25: **while** loop ends

loop $i = 3$

3 † 25: **while** loop ends

loop $i = 4$

4 † 25: **while** loop ends

loop $i = 5$

5 | 25: increment L_5 and replace n by 5

5 | 5: increment L_5 and replace n by 1

Example

 $n = 100$ $L = (0, 2, 0, 0, 2, 0, 0, 0, 0, 0)$ loop $i = 2$ $2 \mid 100$, so increment L_2 and replace n by 50 $2 \mid 50$, so increment L_2 and replace n by 25 $2 \nmid 25$: **while** loop endsloop $i = 3$ $3 \nmid 25$: **while** loop endsloop $i = 4$ $4 \nmid 25$: **while** loop endsloop $i = 5$ $5 \mid 25$: increment L_5 and replace n by 5 $5 \mid 5$: increment L_5 and replace n by 1

...

return $(0, 2, 0, 0, 2, 0, 0, 0, 0, 0)$

What did we get?

$(0, 2, 0, 0, 2, 0, 0, 0, 0, 0)$ tells us

$$100 = 2^2 \times 5^2$$

From there, we can determine $\varphi(100)$ by passing through the loop.

OK, but... lists?

How do we track a *list* of numbers?

Use an **array**, a block of memory.

- need list of 25 int's? `int A[25];`
- need n int's, but don't know n ? `int A[n];`
 - compile w/ `-std=c++11`
- To initialize array to 0, declare instead
`int A[n] { 0 };`

OK, but... lists?

How do we track a *list* of numbers?

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- need list of 25 `int`'s? `int A[25];`
- need n `int`'s, but don't know n ? `int A[n];`
 - compile w/ `-std=c++11`
- To initialize array to 0, declare instead
`int A[n] { 0 };`

Now ready to implement!

Interface

Place in a new directory, factoring

Listing 1: factoring.hpp

```
#ifndef __FACTORING_HPP_  
#define __FACTORING_HPP_  
  
/**  
    A basic factoring algorithm:  
    iterate from 2 to sqrt(n).  
    @param n the number to factor  
    @param primes an array to contain the primes  
    @warning initialize the elements of @c primes  
            to zero!  
*/  
void factors(long n, long * primes);  
  
#endif
```


Implementation

Place in same directory

Listing 2: factoring.cpp

```
#include "factoring.hpp"

void factors(long m, long * primes) {
    long n = m;
    for (long i = 2; n != 1 and i <= m/2; ++i)
    {
        while (n % i == 0) {
            primes[i] += 1;
            n /= i;
        }
    }
}
```

Test program

Place in same directory

Listing 3: test_factoring.cpp (p. 1)

```
#include <iostream>
using std::cin; using std::cout;
using std::endl;

#include "factoring.hpp"

int main() {
    long n;

    cout << "Enter a number to factor --> ";
    cin >> n;

    long m = n/2 + 1;
    long primes[m];
```

Test program

Place in same directory

Listing 4: test_factoring.cpp (p. 2)

```
for (long i = 0; i < m; ++i)
    primes[i] = 0;

factors(n, primes);

for (long i = 0; i < m; ++i) {
    if (primes[i] != 0)
        cout << i << '^' << primes[i] << ' ';
}
cout << endl;
}
```

Compile, execute test

```
$ g++ -c factoring.cpp
$ g++ -o test_factoring -std=c++11 -lm \
    factoring.o test_factoring.cpp
$ ./test_factoring
Enter a number to factor --> 100
2^2 5^2
$ ./test_factoring
Enter a number to factor --> 10
2^1 5^1
```

Things to watch out for

Big-time no-no

Don't forget `-std=c++11`

```
$ g++ -o test_factoring -lm factoring.o \  
    test_factoring.cpp  
test_factoring.cpp: In function 'int main()':  
test_factoring.cpp:14:18: warning: extended initializer  
lists only available with -std=c++11 or -std=gnu++11  
    long primes[m] { 0 };
```

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```

Big-time no-no

Don't forget `-lm`

(You may get away with that last one.)

Homework

p. 89 #5.6, 5.7, 5.10

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This algorithm is not especially efficient

Why not?

- finds n 's prime factorization
- tests divisibility by non-primes
- we could do better if we start with primes

But...

How do we find primes?

Sieve of Eratosthenes

```
given  $n$   
let  $L$  be a list of  $n$  booleans  
    (initialize all  $L$  to True)  
set  $L_1$  to False  
for each  $i \in \{2, \dots, \sqrt{n}\}$   
    if  $L_i$  is True  
        let  $a = i$   
        while  $a < n$   
            add  $i$  to  $a$   
            set  $L_a$  to False  
return  $L$ 
```

Sieve of Eratosthenes

```

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  let  $L$  be a list of  $n$  booleans
    (initialize all  $L$  to True)
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  for each  $i \in \{2, \dots, \sqrt{n}\}$ 
    if  $L_i$  is True
      let  $a = i$ 
      while  $a < n$ 
        add  $i$  to  $a$ 
        set  $L_a$  to False
  return  $L$ 
  
```

Example

$n = 20$

beginning of loop

F			

Sieve of Eratosthenes

given n
 let L be a list of n booleans
 (initialize all L to *True*)
 set L_1 to *False*
for each $i \in \{2, \dots, \sqrt{n}\}$
 if L_i is *True*
 let $a = i$
 while $a < n$
 add i to a
 set L_a to *False*
return L

Example

$n = 20$

$i = 2$ — mark out multiples of 2

F	2		F
	F		F
	F		F
	F		F
	F		F

Sieve of Eratosthenes

given n
 let L be a list of n booleans
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 let $a = i$
 while $a < n$
 add i to a
 set L_a to *False*
return L

Example

$n = 20$

$i = 3$ — mark out multiples of 3

F	2	3	F
	F		F
F	F		F
	F	F	F
	F		F

Sieve of Eratosthenes

given n
 let L be a list of n booleans
 (initialize all L to *True*)
 set L_1 to *False*
for each $i \in \{2, \dots, \sqrt{n}\}$
 if L_i is *True*
 let $a = i$
 while $a < n$
 add i to a
 set L_a to *False*
return L

Example

 $n = 20$ $i = 4 - L_4 = F$: not prime; **skip!**

F	2	3	F
	F		F
F	F		F
	F	F	F
	F		F

Sieve of Eratosthenes

given n
let L be a list of n booleans
(initialize all L to *True*)
set L_1 to *False*
for each $i \in \{2, \dots, \sqrt{n}\}$
 if L_i is *True*
 let $a = i$
 while $a < n$
 add i to a
 set L_a to *False*
return L

Example

 $n = 20$ $i = 5 > \sqrt{20}$ – end loop, true
entries prime!

F	2	3	F
5	F	7	F
F	F	11	F
13	F	F	F
17	F	19	F

Observation

Theorem

Factoring n requires us to test at most $\sqrt[4]{n}$ numbers for primality.

Proof.

- Prime factor of n must be smaller than \sqrt{n} .
 - Sieve of Eratosthenes needs \sqrt{m} tests to find all primes less than m .
- \therefore Need $\sqrt{\sqrt{n}} = \sqrt[4]{n}$ tests.



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- \therefore Need $\sqrt{\sqrt{n}} = \sqrt[4]{n}$ tests.



Time to implement the sieve!

Creating lists on-the-fly

Two kinds of array allocation

static create array using `Type var[num];`

dynamic create array using

`Type * var = new Type[num];`

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Why two?

- static allocation...
 - *more efficient* if num is known constant (e.g., “25”)
 - *unsafe* if data needed outside function
 - memory **will be** trashed!
- dynamic allocation
 - when finished w/array, requires `delete [] var;`
 - *safe* to pass outside function

Creating lists on-the-fly

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 - memory **will be** trashed!
- dynamic allocation
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sieve's list of primes: *dynamic*

- don't know how many
- need to return to caller

Interface

Place in factoring directory

Listing 5: sieve.hpp

```
#ifndef __SIEVE_H_  
#define __SIEVE_H_  
  
/**  
    Sieve of Eratosthenes:  
    generate table of primes.  
    @param n find primes <= n  
    @param primes array of primes  
    @return number of primes found */  
long sieve(long n, long * & primes);  
  
#endif
```

Implementation

In same directory

Listing 6: sieve.cpp (p. 1)

```
#include <cmath>
using std::sqrt;

#include "sieve.hpp"

long sieve(long n, long * & primes) {
    long m = long(sqrt(n));
    primes = new long[n];

    long num_primes = 0;
    bool * theSieve = new bool[n];
    for (long i = 2; i < n; ++i)
        theSieve[i] = true;
```

Implementation

In same directory

Listing 7: sieve.cpp (p. 2)

```
long i;
for (i = 2; i < m + 1; ++i) {
    if (theSieve[i] == true) {
        primes[num_primes] = i;
        ++num_primes;
        long a = i;
        while (a < n) {
            theSieve[a] = false;
            a += i;
        }
    }
}
```

Implementation

In same directory

Listing 8: sieve.cpp (p. 3)

```
for (/* */ ; i < n ; ++i) {
    if (theSieve[i] == true) {
        primes[num_primes] = i;
        ++num_primes;
    }
}
delete [] theSieve;
return num_primes;
}
```


Test program

Place in same directory

Listing 9: test_sieve.cpp (p. 1)

```
#include <iostream>
using std::cin; using std::cout;
using std::endl;

#include "sieve.hpp"

int main() {
    long * primes;
    long n;

    cout << "This program finds all primes ";
    cout << "less than your choice of number.\n";
    cout << "Please choose a number --> ";
    cin >> n;
```

Test program

Place in same directory

Listing 10: test_sieve.cpp (p. 2)

```
long np = sieve(n, primes);  
cout << "There are " << np << " primes:\n";  
for (long i = 0; i < np; ++i)  
    cout << primes[i] << ", ";  
cout << endl;  
delete [] primes;  
}
```

Prime pairs, yet
again

Euler's totient
function

A first factoring
algorithm

**A better factoring
algorithm**

Back to the totient

Counting pairs

Answering the
question

Summary

p. 90 #5.11

Homework

Outline

- 1 Prime pairs, yet again
- 2 Euler's totient function
 - A first factoring algorithm
 - A better factoring algorithm
 - Back to the totient
- 3 Counting pairs
- 4 Answering the question
- 5 Summary

Recall our algorithm for p_n

```
given  $n$   
let  $s = 0$   
for each  $k \in \{2, \dots, n\}$   
    add  $\varphi(k)$  to  $s$   
multiply 2 to  $s$   
add 1 to  $s$   
divide  $s$  by  $n^2$   
return  $s$ 
```

...we still need to compute the totient

Interface

Place in factoring folder

Listing 11: totient.hpp

```
#ifndef __TOTIENT_HPP__  
#define __TOTIENT_HPP__  
  
long totient(long n, long * primes);  
  
#endif
```

One more property

Theorem

$$\varphi(n) = n \times \left(\frac{p_1 - 1}{p_1}\right) \times \left(\frac{p_2 - 1}{p_2}\right) \times \cdots \times \left(\frac{p_\ell - 1}{p_\ell}\right)$$

where p_1, p_2, \dots, p_ℓ are the prime factors of n .

Proof.

Think about it a moment...

One more property

Theorem

$$\varphi(n) = n \times \left(\frac{p_1 - 1}{p_1}\right) \times \left(\frac{p_2 - 1}{p_2}\right) \times \cdots \times \left(\frac{p_\ell - 1}{p_\ell}\right)$$

where p_1, p_2, \dots, p_ℓ are the prime factors of n .

Proof.

Recall that

$$\varphi(n) = (p_1^{k_1} - p_1^{k_1-1}) \cdots (p_\ell^{k_\ell} - p_\ell^{k_\ell-1}).$$

Factor each factor's common factor:

$$\varphi(n) = (p_1^{k_1-1} \cdots p_\ell^{k_\ell-1}) \times [(p_1 - 1) \cdots (p_\ell - 1)].$$

The leftmost product can be rewritten as

$$\varphi(n) = \frac{n}{p_1 \cdots p_\ell} \times [(p_1 - 1) \cdots (p_\ell - 1)],$$

and we are done.

Implementation

Place in factoring folder

Listing 12: totient.cpp

```
#include <cmath>
using std::sqrt;

#include "totient.hpp"

long totient(long n, long * primes) {
    if (n < 0) return 0;

    long result = n;
    for (long i = 0; n != 1 and primes[i] <= n; ++i) {
        if (n % primes[i] == 0) {
            result /= primes[i];
            result *= primes[i] - 1;
        }
    }
    return result;
}
```

Test program

Place in factoring folder

Listing 13: test_totient.cpp

```
#include <iostream>
using std::cin; using std::cout;
using std::endl;

#include "sieve.hpp"
#include "totient.hpp"

int main() {
    long n;
    cout << "This programing computes ";
    cout << "the totient of an integer.\n";
    cout << "Please input a number --> ";
    cin >> n;
    long *primes;
    long np = sieve(n, primes);
    cout << totient(n, primes) << endl;
    delete [] primes;
}
```

Compile, execute

```
$ g++ -c sieve.cpp totient.cpp
$ g++ -o test_totient sieve.o totient.o \
    test_totient.cpp
$ ./test_totient
Please input a number --> 100
40
```

Outline

- 1 Prime pairs, yet again
- 2 Euler's totient function
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 - A better factoring algorithm
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- 4 Answering the question
- 5 Summary

Recall our algorithm for p_n

```
given  $n$   
let  $s = 0$   
for each  $k \in \{2, \dots, n\}$   
    add  $\varphi(k)$  to  $s$   
multiply 2 to  $s$   
add 1 to  $s$   
divide  $s$  by  $n^2$   
return  $s$ 
```

Implementation

Place in `relprime_pairs` folder

Listing 14: `totient_pairs.cpp` (p. 1)

```
#include <iostream>
using std::cin; using std::cout;
using std::endl;
#include <iomanip>
using std::setprecision;

#include "../factoring/sieve.hpp"
#include "../factoring/totient.hpp"

/**
    Calculates probability that two int's
    chosen in {1,2,...,n} are rel prime,
    up to  $n=10^6$ .
*/
```

Implementation

Place in `relprime_pairs` folder

Listing 15: `totient_pairs.cpp` (p. 2)

```
int main() {  
    const long N = 10000000;  
    const long UPDATE = 100000;  
  
    long * primes;  
    long np = sieve(N, primes);  
  
    long count = 0;  
    cout << setprecision(20);
```

Implementation

Place in relprime_pairs folder

Listing 16: totient_pairs.cpp (p. 3)

```
cout << setprecision(20);
for (long k = 1; k <= N; ++k) {
    count += totient(k, primes);
    if (k % UPDATE == 0) {
        cout << k/1000 << " thousand \t";
        cout << double(2*count - 1)
            / double(k*k) << endl;
    }
}
delete [] primes;
return 0;
}
```


Compiling, executing

Prime pairs, yet
again

Euler's totient
function

A first factoring
algorithm

A better factoring
algorithm

Back to the totient

Counting pairs

Answering the
question

Summary

```
$ g++ -Ofast -o totient_pairs \  
    totient_pairs.cpp \  
    ../factoring/sieve.o \  
    ../factoring/totient.o  
$ ./totient_pairs  
100 thousand      0.60793015070000000488  
200 thousand      0.607929945875000044  
300 thousand      0.60792774407777783185  
...  
800 thousand      0.60792796007343752329  
900 thousand      0.60792736490740739708  
1000 thousand     0.60792710478300004961
```

Outline

- 1 Prime pairs, yet again
- 2 Euler's totient function
 - A first factoring algorithm
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- 5 Summary

The “point”

Question

So what is $\lim_{n \rightarrow \infty} p_n$?

We found

$$\lim_{n \rightarrow \infty} p_n \approx 0.607927.$$

Online Encyclopedia of Integer Sequences: $6/\pi^2!$

The “point”

Question

So what is $\lim_{n \rightarrow \infty} p_n$?

We found

$$\lim_{n \rightarrow \infty} p_n \approx 0.607927.$$

Online Encyclopedia of Integer Sequences: $6/\pi^2!$

???

How does π figure into this?

Basel problem (famous)

Find

$$\sum_{n=1}^{\infty} \frac{1}{n^2}.$$

How does π figure into this?

Basel problem (famous)

Find

$$\sum_{n=1}^{\infty} \frac{1}{n^2}.$$

Solution: (Euler, 1734)

$$\sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{\pi^2}{6}.$$

Proof.

Hard. Go bother Dr. Hornor or Dr. Ding or Dr. Kohl.



That doesn't explain jack squat.

Look at

$$\begin{aligned} \sum_{n=1}^{\infty} \frac{1}{n^2} &= \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{4^2} + \frac{1}{5^2} + \frac{1}{6^2} + \dots \\ &= \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{2^4} + \frac{1}{5^2} + \frac{1}{2^2 \times 3^2} + \dots \\ &= \prod_p \left(1 + \frac{1}{p^2} + \frac{1}{p^4} + \frac{1}{p^6} + \dots \right) \\ &= \prod_p \left(\frac{1}{1 - \frac{1}{p^2}} \right). \end{aligned}$$

So

$$\therefore \prod_p \left(1 - \frac{1}{p^2} \right) = \frac{1}{\sum 1/n^2} = \frac{6}{\pi^2}.$$

...and?

Probability two integers:

- divisible by 2: $\frac{1}{2^2}$, so not: $1 - \frac{1}{2^2}$
- divisible by 3: $\frac{1}{3^2}$, so not: $1 - \frac{1}{3^2}$
- divisible by 5: $\frac{1}{5^2}$, so not: $1 - \frac{1}{5^2}$

...

SO...

...and?

Probability two integers:

- divisible by 2: $\frac{1}{2^2}$, so not: $1 - \frac{1}{2^2}$
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- divisible by 5: $\frac{1}{5^2}$, so not: $1 - \frac{1}{5^2}$

...

SO...

Probability two large integers relatively prime:

$$\prod_{p < \text{larger}} \left(1 - \frac{1}{p^2}\right)$$

...and?

Probability two integers:

- divisible by 2: $\frac{1}{2^2}$, so not: $1 - \frac{1}{2^2}$
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- divisible by 5: $\frac{1}{5^2}$, so not: $1 - \frac{1}{5^2}$

...

SO...

Probability two large integers relatively prime:

$$\prod_{p < \text{larger}} \left(1 - \frac{1}{p^2}\right)$$

Thus

$$\lim_{n \rightarrow \infty} p_n = \prod_p \left(1 - \frac{1}{p^2}\right) = \frac{6}{\pi^2}.$$

Prime pairs, yet
again

Euler's totient
function

A first factoring
algorithm

A better factoring
algorithm

Back to the totient

Counting pairs

Answering the
question

Summary

p. 90 #5.12

Homework

Outline

- 1 Prime pairs, yet again
- 2 Euler's totient function
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 - A better factoring algorithm
 - Back to the totient
- 3 Counting pairs
- 4 Answering the question
- 5 Summary

Summary

- Math stuff
 - Euler totient function, properties
 - Sieve of Eratosthenes
 - Online Encyclopedia of Integer Sequences
 - π turns up in the strangest places!
- Programming stuff
 - arrays
 - static array creation
 - dynamic array creation
 - new and delete []